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W. S.

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U.S. Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number. **Application Number** 10/564,568 Filing Date 13 January 2006 TRANSMITTAL First Named Inventor Sergey Nikolaevich Zheltov **FORM** Art Unit Not yet assigned **Examiner Name** Not yet assigned (to be used for all correspondence after initial filing) Attorney Docket Number 42390P16120 Total Number of Pages in This Submission **ENCLOSURES** (Check all that apply) After Allowance Communication to TC Fee Transmittal Form Drawing(s) Appeal Communication to Board Licensing-related Papers Fee Attached of Appeals and Interferences Appeal Communication to TC Petition Amendment/Reply (Appeal Notice, Brief, Reply Brief) Petition to Convert to a **Proprietary Information** After Final **Provisional Application** Power of Attorney, Revocation Status Letter Affidavits/declaration(s) Change of Correspondence Address Other Enclosure(s) (please Identify Terminal Disclaimer Extension of Time Request below): Return Receipt Postcard Request for Refund **Express Abandonment Request** CD, Number of CD(s) Information Disclosure Statement Landscape Table on CD Certified Copy of Priority Remarks Express Mail No. EV 841 071 808 US Document(s) Reply to Missing Parts/ Incomplete Application Reply to Missing Parts under 37 CFR 1.52 or 1.53 SIGNATURE OF APPLICANT, ATTORNEY, OR AGENT BLAKELY, SOKOLOFF, TAYLOR & ZAFMAN, LLP Firm Name 12400 Wilshire Boulevard Seventh Floor, Los Angeles, CA 90025-1030 Signature Printed name James C. Scheller Date Reg. No. 31.195 2006 **CERTIFICATE OF MAILING** I hereby certify that this correspondence is being deposited with the United States Postal Service on the date shown below with sufficient postage as first class mail in an envelope addressed to: Mail Stop-RCT, Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.

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Jessica Savage

Signature

Typed or printed name

РОСПАТЕНТ

Федеральное государственное учреждение «Федеральный институт промышленной собственности Федеральной службы по интеллектуальной собственности, патентам и товарным знакам» (ФГУ ФИПС)

Бережковская наб., 30, корп. 1, Москва, Г-59, ГСП-5, 123995 Телефон 240- 60- 15. Телекс 114818 ПДЧ. Факс 234- 30- 58

Наш № 20/12-260

«11» мая 2006 г.

СПРАВКА

Федеральный институт промышленной собственности (далее - Институт) настоящим удостоверяет, что приложенные материалы являются точным воспроизведением первоначального заявления, описания, формулы, реферата и чертежей (если имеются) международной заявки № РСТ/RU2003/00306, поданной в Институт как в Получающее ведомство в соответствии с Договором о патентной кооперации 15 июля 2003 года (15.07.2003).

CERTIFIED COPY OF PRIORITY DOCUMENT

И.О. заведующего отделом 20

Т.Ф.Владимирова

Копия для получающего ведомства

PCT

ЗАЯВЛЕНИЕ

Нижеподписавшийся просит рассматривать настоящую международную заявку в соответствии с Договором о патентной кооперации

Заполняется получающим ведомством — PCT/RU 0 3 / 0 0 3 0 6						
Номер международной заявки						
15 ИЮЛЯ 2003 (15.07.2003). Дата международной подачи						
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урафа 1	monitoring for symmet				
Графа II	ЗАЯВИТЕЛЬ Г	. Ланное г	INIIO SRIISPTOS T	акже изобретателем	·
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	5252, Moscow, Sokol-10 vsky per., 14	Business	s Center,	Телепринтер №	
				Регистрационный №	•
				заявителя в Ведомстве	. ,
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заявителем дл	ія: Государств	государств, кр	оме США	дополнительной	
Графа III	ДРУГИЕ ЗАЯВИТЕЛИ И/ИЛИ	І (ДРУГИЕ)	ИЗОБРЕТАТЕ	ЛИ	
	Фамилия указывается перед именем, для юридическ и включать почтовый индекс и название страны. Ес			Данное лицо является:	
	изано, то таковым будет считаться страна указа			только заявителем:	
ZHELTO	N Sergei Nikolaevich			Х заявителем и изобретателем	4
RU, 60	3950, Nizhny Novgorod,	Turgene	v st.,30		
			i	только изобретателем (если с этот бокс, то ниже заполня	тмечен ть
				не требуется)	
				Регистрационный № заявителя в Ведомстве	
				заявителя в Ведометве	
осударство (п	т.е. страна) гражданства: RU	Гос	сударство(т.е. сп	прана) местожительства: RU	
јанное лицо я		всех указанных		голько США государств, указан	ных в
аявителем для	я: государств г	осударств, кро	ome CILIA X	дополнительной г	рафе
х Друг	гие заявители и/или (другие) изобре [.]	гатели назван	на листе прод	должения	
Графа IV	АГЕНТ ИЛИ ОБЩИЙ ПРЕДСТ	ГАВИТЕЛЬ;	или адрес,	для переписки	
казанное ниж	ке лицо настоящим назначается (назначе	ено) представл	ATR -	агента общего	
интересы заяв	вителя(ей) в компетентных международн	ных органах в н	качестве:	представите	еля
	Фамилия указывается перед именем, для юрч Адрес должен включать почтовый индекс и н	· ·	· · · · · ·	Телефон № 925-16-61	
OBSCHE	ESTVO S OGRANICHENNOI			Телефакс №	
OTVETS	STVENNOSTJU "SOJUZPATEI			924-95-40	
RU, 10	03735, Moscow, ul. Ili:	inka, d.	5/2	Телепринтер №	
	•				ĺ
			ì	Регистрационный №	
			·	агента в Ведомстве	
	с для переписки: Пометить этот бокс, с				
указа	нный выше адрес используется только в	ак специальнь	ій адрес для пере	писки	

	Графа III ДРУГИЕ ЗАЯВИТЕЛИ И/ИЛИ (ДРУГИЕ) ИЗОБРЕТАТЕЛИ Если ни одна из следующих подграф не используется, этот лист не включается в заявление						
·	Имя и адрес: (Фачилия указывается перед вменем, для юридического лица - полное уставное наиме пис. Адрес дляжен включать почтовый индекс в название страны. Если госудаются о местолуство ст	ново: Данное лицо является:					
	ВRATANOV Stanislav Viktorovich	только заявителем: х заявителем и изобретателем					
•	RU, 603950, Nizhny Novgorod, Turgenev st., 30	только изобретателем (если стмече этот бокс, то ниже заполныть не требуется)					
		Регистрационный № заявителя в Ведомстве					
	Государство (т.е. страна) гражданства: RU Государство(т.е. страна) местожительства: RU					
	Данное лицо является всех указанных всех указанных государств, кроме США	х только США государств, указанных в дополнительной графе					
- 14	Имя и адрес: (Фачилия указынается перед именем, для юридического лица - полное уставное наимен пис. Адрес дольжен включать почтовый индекс и название страны. Если государство местожительство янизу не будет указанного в данной графе адреса) BELENOV Roman Alexeevich RU, 603950, Nizhny Novgorod, Turgenev st., 30	только заявителем: заявителем и изобретателем только изобретателем (если отымечен					
		этот бокс, то ниже заполнянь не требуется) Регистрационный №					
	осударство <i>(т.е. страна)</i> гражданства: Государство <i>те</i>	заявителя в Ведомстве					
\perp	RU всех указанных всех указанных посударство м.е	. <i>страна)</i> местожительства:					
32	аявителем для: государств государств, кроме США х	дополнительной графе					
1""	мя и адрес: (Фамилия указывается перед именем, для юридического лица - полное уставное наименой се. Адрес должен включать почтовый индекс и название страны. Если государство местожительства изу не будет указанного в данной графе адреса) KNYAZEV Alexandr Nikolaevich RU, 603950, Nizhny Novgorod, Turgenev st., 30	Данное лицо является: только заявителем: заявителем и изобретателем только изобретателем (если отмечен этот бокс, то ниже заполнять не требуется)					
		Регистрационный № заявителя в Ведомстве					
		страна) местожительства: RU					
385	всех указанных государств, кроме США	только США государств, указанных в дополнительной графе					
нис	19 И адрес:(Фамилия указывается перед именем, для юридического лица - полное устанное наименове . Адрес далжен включать почтовый индекс и название страны. Если государство местожительства зу не будет указани, то таковым будет считоться страна указанного в данной графе адреса)	только заявителем: заявителем и изобретателем только изобретателем (если отмечен этот бокс, то ниже заполнять не требуется)					
-00	Whomewo (ii)	Регистрационный № заявителя в Ведомстве					
	Wilder Private - Transport	страна) местожительства:					
	всех указанных всех указанных государств, кроме США	только США государств, указанных в дополнительной графе					
[Другие заявители и/или (другие) изобретатели названы на другом ли	сте для продолжения					

									
	аф:				·			быть	отмечен как минимум офін бокс
Настоящим делаются следующие указания в соответствии с правилом 4.9(а):									
Региональный патент							•		
	 АР Патент ARIPO: GH Гана, GM Гамбия, KE Кения, LS Лесото, MW Малави, MZ Мозамбик, SD Судан, SL Сьерра- Леоне, SZ Свазиленд, TZ Объединенная Республика Танзания, UG Уганда, ZH Замбия, ZW Зимбабвс, а также любое другое государство, являющееся Договаривающимся государством Протокола Хараре и РСТ(если испрашивается иной вид охраны или статус, написать на пунктирной линии): ЕА Евразийский патент: АМ Армения, AZ Азербайджан, BY Беларусь, KG Кыргызстан, KZ Казахстан, 								
	MD Республика Молдова, RU Российская Федерация, ТЈ Таджикистан, ТМ Туркменистан, а также любое другое государство, являющееся Договаривающимся государством Евразийской патентной конвенции и РСТ								
(X	EP	Европейский патент: АТ Австри DK Дания, ES Испания, FI Финля LU Люксембург, MC Монако, NL государство, являющееся Договар	нди: Ни	я, FR церла	1 Франция, GB Великобритания, анды, РТ Португалия, SE Швеци	, GR ія, Т	Грец R Туј	ия, ІЕ Ирландия, ІТ Италея,
	ОА Патент ОАРІ: В Буркина Фасо, В Бенин, С Г Центральная Африканская республика, С С Конго, С І Кот д'Ивуар, С М Камерун, G А Габон, G N Г Винея, G Q Экваториальная Г Винея, G W Г Винея-Бисау, M L Мали, MR Мавритания, NE Нигер, S N Сенегал, T D Чад, T G Того а также любое другое государство, являющееся членом ОАРІ и Договаривающимся государством Р С Т (если испрашивается иной вид охраны или статус, написать на пунктирной линии): Национальный патент (если испрашивается иной вид охраны или статус, написать на пунктирной линии):								
				,					
			ъединенные Арабские Эмираты	N N		Гамбия	X	••	I Оман Новая Зеландия
[3]			тигуа и Барбуда бания	図		Хорватия	$\overline{\mathbf{z}}$		новая зеландия Филиппины
N N	AL	A AUIC A A	рмения	$\overline{\mathbf{x}}$		ВенгрияИндонезия	X		Польша
			стрия	$\overline{\mathbf{x}}$		Израиль		PT	Португалия
X	AT	AR	стралия	X		Индия:	X	RO	Румыния
Ä			вербайджан		is	Исландия	\boxtimes		Российская Федерация
$\overline{\mathbf{Z}}$			сния и Герцеговина	$\hat{\mathbf{x}}$:		·	•
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\square			ларусь	-		тическая республика	四 -	-SI	Словения
			лиз	\mathbf{Z}	KR	Республика Корея	図	SK	Словакия
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\mathbf{x}			LI Швейцария и Лихтенштейн	図	LC	Сент-Люсия	図	TJ	Таджикистан
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A			миника .	X		Республика Молдова	×		Соединенные Штаты Америки
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			задор	X	MK	Бывшая Югославская респуб-			Узбекистан
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	ксы,	заре	зервированные для указания госу.	даро	·				уска данного листа
	•••••	• • • • • • • • • • • • • • • • • • • •	•••••			······································		•••••	
пр пр ука до:	авил ивед азан лжн	юм ценні ия г счі	ние о предварительных указа 4.9(b), делает также все указан ого в Дополнительной графе в к подлежат подтверждению, и что итаться изъятым заявителем на м ставлено в получающее ведомств	ия, ачес люб юме	допу тве и бое у нт и	стимые в соответствии с РСТ, есключенных из данного упоминиказание, не подтвержденное до стечения этого срока. (Подтверс	за нани: ист	искл я, и за ечени	ночением указания (указаний), аявляет, что эти дополнительные

	ление на приоритет	<		· · · · · · · · · · · · · · · · · · ·	
Настояшим заявляетс	я прноритет следующей предше	ствующей заявки(ок):		· · · · ·	
Дата подачи	Номер	Если п	IDEAUIECTRYIOIUAR 320RK2 AR	7007.00	
предшествуюшей зая: (день/месяцігод)		национальной заявкой:	редшествующая заявка является: региональной заявкой: международно		
(1)		страна .	региональное ведомство	международнов заявк получающее веломет	
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5)					
Поправиления				•	
——————————————————————————————————————	аявления на приоритет указаны гву поручается подготовить и на в том случае, если предшествую				
Торговой Организац рафа VII МЕЖДУ Выбор международного более международного КИ SA / RU Просьба об использова	ощей заявкой является заявка кой конвенции по охране прог ции, в которую была подана р УНАРОДНЫЙ ПОИСКОВЬ по поискового органа (ISA) (есл ных поисковых органа, указать в нии результатов ранее проведи ародного поискового органа ран Номер	одиняя заявка (правило БІЙ ОРГАН и компетентными в провыбранный поисковый оро	ости или одна страна- о 4.10(b)(ii) ведении международного ган; можно использовать	поиска являются два двубуквенный код): к был уже проведен	
рафа VIII ДЕКЛ	АРАЦИИ				
анное заявление содер	жит следующие декларации (ни	же отметить		Vonum	
еобходимые боксы и ук еклараций):	сазать в правой колонке количес	тво каждого типа		Количество деклараций	
Графа VIII (i)	Декларация об удостоверении л	ичности изобретателя			
Графа VIII (ii)	Декларация о правомочности за подачи подавать заявку и получ	явителя на дату межлуна	ародной		
Графа VIII (iii)	Декларация о правомочности за подачи на заявление о приорите заявление предшест	явителя на дату междуна те в случае, если он не я	: продной вляется		
Графа VIII (iv)	Декларация об авторстве на изоб Соединенных Штатов Америки	•	ния :		
Графа VIII (v) Д	декларация о не наносящих ушер отсутствия новизны	об раскрытиях или изъят	иях из-за		

Графа IX КОНТРОЛЬНЫЙ ПІ	ЕРЕЧЕНЬ; ЯЗ	ык подачи	
Настоящая международная заявка сод	ержит:	CREDITORIUS PORTE COMPANIA	Кол-во
(а) следующее количество листов на	•	COOMCOMO MICHIGAN CONTRACTOR CONT	прило
бумажном носнтеле:	_	соответствующие боксы и указать с правса количество приложений каждого вида):	жений
заявление(включая декларации)	: 5	количество приложении кажоого вида).	
описание (исключая перечень	15	. 1. П лист расчета пошлин	:
последовательностей)	7	2. П оригинал отдельной доверенности	
формула	•	2. Оригинал отдельной доверенности	:
реферат	: 1	3. оригинал генеральной доверенности	:
чертежи	: 3	4. П копия генеральной доверенности; ссылка	
Предварительное число листов	: 31		
	. 21	на номер, если имеется:	:
часть описания с перечнем после-		5. разъяснения по поводу отсутствия подписи	:
довательностей (действительное		6. 🔲 приоритетный(ые) документ(ы), указанный	
число листов, представленных на	i	в графе VI под №	:
бумажном носителе, независимо		7. Перевод международной заявки на	
от представления в машиночитаем	10ŭ	(язык)	:
форме; см. ниже пункт (b)	:	8. информация о депонировании микроорганизмов	
•		или другого биологического материала	:
Общее число листов	: 31		
•		9. перечень последовательностей в машиночитаемой	
(b) перечень последовательностей пре	едставлен в	форме(указать тип и число носителей (дискета,	
машиночитаемой форме	.]	CD-ROM, CD-R или иное))	
(і) П только (в соответствии с разд	16110M 801(a)(i))	(і) копия, представленная для целей международного	
(1) TOJIBKO (B COOTBETETBUR C PASS	(C)10M 001(a)(1))	поиска в соответствии с правилом 13 ter (и не	
(ii) как приложение к представле	нному на	являющаяся частью международной заявки) :	
бумажном носителе(в соотве	тствии с	(ii) (только в случае, если слева отмечены бокс(b)(i)	
разделом 801(а)(іі))		или (b)(ii)) дополнительно представленная копия,	
	i	если допустимо, копия для целей международного	
Тип и количество носителей (дискет	a, CD-ROM,	поиска в соответствии с правилом13 ter	
CD-R или другое), на котором предста	авлен перечень	(iii) вместе с соответствующим представлением	
последовательностей (дополнительно	к указанному в	перечня последовательностей, как его заявление	•
пункте 9(іі) в правой колонке):		отмечено слева :	
		10. Ц иное (указать)	
Фигура чертежей, предлагаемая		Язык подачи	
для публикации с рефератом:		международной заявки: GB	
Графа Х ПОДПИСЬ ЗАЯВИТЕ	ЕЛЯ, АГЕНТА	ИЛИ ОБЩЕГО ПРЕДСТАВИТЕЛЯ	
	чь каждого подпі	исавшего и указать, в каком качестве он подписал заявление (если это не очес	видно
из данных, приведенных в заявлении).		·	
		General Director of 000 "Sojuzpatent"	
		Dalitana C. D	
		Felitsyna/S.B.	
	- Заполняется	я получающим ведомством	
. Дата фактического получения	T F TATA TOTAL	2. Чертежи	:
международной заявки:	RRONET	2003 (15.07.2003)	
 Исправленная дата при более поздне получении страниц или чертежей, д предполагаемую международную за 	оукомплектовыв	зающих	
. Дата своевременного получения тре		не получ	нены:
•	-		ļ
исправлений согласно статье 11(2) Р	C1.		
Международный поисковый орган (если компетентны два и более):	ISA/RU 6	Направление копии для поиска задержано впредь до уплаты пошлины за поиск	
	3200000000	a May muona nu M fiona	
Дата получения регистрационного экз Международным бюро:		я Международным бюро	

A Method of Efficient Performance Monitoring for Symmetric Multi-Threading Systems

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BACKGROUND

1. FIELD

The present invention relates generally to performance measurement techniques and, more specifically, to measurement of performance of an execution thread within a symmetric multi-threading (SMT) system.

2. DESCRIPTION

It is a general practice to increase the computational performance by organizing parallel program execution. There are a number of methods to achieve this, including, but not limited to, out-of-order instruction execution, multiple data operands, shared memory multi-processor systems, distributed computations, and so forth. One of the popular and relatively inexpensive approaches is to combine multiple execution cores within one physical processor, or even provide separate execution state containers and control logic to share multiple processing units of a physical processor. The latter statement is applicable, for example, to the Hyper-Threading technology commercially available from Intel Corporation, which provides better utilization of various execution units incorporated in a processor.

Measurement of a processor's (program's) performance is one of the main tasks to be solved when building an efficient computational system. For single processor systems, performance monitoring is a matter of correctly written software, given that the processor (or other hardware components) provides the necessary resources. The performance monitoring task may be more difficult for SMT systems: performance monitoring

hardware support may vary considerably, and the interaction between hardware and software parts of performance monitoring system becomes more complicated.

Possible difficulties that can arise include the lack of performance monitoring resources (e.g., performance counters) to monitor the activity of all processing units (e.g., logical threads or processors) within a physical package, and no hardware support of asynchronous and independent measurements performed on a per-thread (per-logical processor/unit) basis.

Therefore, a need exists for the capability to efficiently monitor the performance of multi-threading systems taking into account the possible lack of hardware resources.

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BRIEF DESCRIPTION OF THE DRAWINGS

The features and advantages of the present invention will become apparent from the following detailed description of the present invention in which:

Figure 1 is a diagram illustrating the dedication of hardware resources to execution threads according to an embodiment of the present invention;

Figure 2 is a flow diagram illustrating the initiation of the performance monitoring process according to an embodiment of the present invention; and

Figure 3 is a flow diagram illustrating the completion of the performance monitoring operation according to an embodiment of the present invention.

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DETAILED DESCRIPTION

Embodiments of the invention described herein may be applicable to performance monitoring conducted on an execution thread basis within a symmetric multi-threading (SMT) system. One embodiment of the present invention may be used in a system built on Intel Corporation's Hyper-Threading (HT) technology to enable effective performance monitoring on a per logical processor basis.

Reference in the specification to "one embodiment" or "an embodiment" of the present invention means that a particular feature, structure or characteristic described in connection with the embodiment is included in at least one embodiment of the present invention. Thus, the appearances of the phrase "in one embodiment" appearing in various places throughout the specification are not necessarily all referring to the same embodiment.

It is not always possible to provide independent hardware support for simultaneous monitoring of multiple execution threads or logical execution modules (e.g., logical processors for HT). Thus, many useful measurements can be performed for either all execution threads or for a specified subset, depending on a particular hardware implementation. Embodiments of the present invention relate to the case of limited performance monitoring resources and enable quasi-independent measurements for each execution thread or logical execution unit. That is, whenever a thread (logical unit) initiates measurements, the overall performance monitoring results are computed correctly, but the distribution of the results to any particular thread (logical unit) depends on a particular hardware implementation.

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The following definitions may be useful in understanding embodiments of the present invention described herein.

A performance monitoring unit is a device (whether external, integrated, or a specific functional block within a primary unit) intended for measuring (monitoring) operational characteristics of a primary device (unit) or system.

An execution thread is a program to be executed by a processing unit (e.g., processor) independently and (if possible) concurrently with other programs, and the state of the processing unit (execution context) associated with such a program.

A logical execution unit is a specific processing unit that executes a program concurrently with other processing units, maintains a program execution state, and shares system resources with similar units within a primary processing unit.

One logical execution unit is supposed to run one execution thread (program) at a time. Therefore, for purposes of describing embodiments of the present invention there is no essential difference between the two terms. The methods described herein may be applicable to any processing system that may have performance monitoring resources shared between multiple processing units as well as multiple program threads as the latter are supported by hardware.

Hereinafter the term 'execution unit' denotes both an execution thread and a logical execution unit.

Figure 1 illustrates the structure of a performance monitoring unit (PMU) and three types of resource sharing that may occur in a symmetric multi-threading system. A PMU comprises counter logic 10, control logic 12, and execution unit indicator logic 14. In some embodiments, the execution unit indicator logic may be a part of the control logic. In

a system that supports multiple execution units (EUs) within a physical package and provides each EU with a separate PMU for any given performance monitoring functionality, all performance monitoring data may be collected independently and asynchronously on an EU's demand. There are, however, a number of systems with limited PMU resources (e.g., Intel Corporation's Pentium4 processor with Hyper-Threading technology enabled) that need to be shared between multiple execution units. One of the examples of such sharing may be a system that has only one PMU that is capable of collecting performance data for either one or all execution units. The former case (one EU to be monitored) generally results in undercounting of performance data, while the latter case will produce overcounted results. To handle both cases, a system that implements the present invention needs to emulate the execution unit indicators 16 for each EU by means of a request allocation as described below.

Most of the current state-of-the-art systems provide a capability to set up a PMU to collect performance data for a subset of execution units by furnishing additional EU-indicators 18. Typically, the number of additional EU-indicators equals the number of execution units within a package, otherwise, if the number of EU-indicators appears to be less, the above described single EU-indicator conditions hold true for this case.

The performance monitoring process is illustrated in Figures 2 and 3.

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According to embodiments of the present invention, the performance monitoring is started or stopped upon a request from an execution unit. A system implementing the present method should be capable of maintaining the correct sequence of such requests, insuring that a stop request always follows a start request or establishing the start/stop correspondence in any other applicable manner, e.g., providing a nested request support or ignoring excessive requests. As the requests may appear simultaneously, a special arbitration step 20 may be used to guarantee the exclusive use of a PMU. Once exclusive execution is acquired, the start request is allocated at block 22, that is, a special table (provided for this purpose) or PMU (if supported by hardware) field may be filled with a value indicating that a request to start performance monitoring operation is pending for a specific execution unit. If there is only one request currently allocated, the PMU may be programmed at block 24 to start collecting performance monitoring data for the EU that allocated the request. At block 26, the PMU counter may be set to a predefined value if the hardware supports counter initialization; otherwise, the current counter value may be stored in a special memory area as an initial value and may be later subtracted from a final

value when the performance monitoring operation stops. In case there is more than one request already allocated, the requesting EU may be added to the set of EUs the PMU currently collects performance data for if such a possibility is supported by the PMU's hardware (PMU has a free EU-indicator).

Thus, performance monitoring process starts, and one counter of one PMU accumulates performance data for all execution units as they request this operation.

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To stop the operation for an EU, a stop request may be issued by this execution unit. The arbitration may be performed at block 100 to acquire exclusive processing of each stop request. Then, at block 102, the stop request may be removed from the special table (see above) or PMU EU-indicator field (if supported by hardware). If there are no more requests allocated, the PMU may be programmed to stop collecting performance data at block 104. The final performance value may then be obtained at block 106. If there are requests from other EUs, active or pending, and the current EU belongs to the set of EUs the PMU collects data for (i.e., the request is active), the PMU may be programmed to stop collecting data for the current EU (if such a possibility is supported by the PMU's hardware). Then, one skilled in the art will recognize the option, based on the knowledge of a particular system architecture and hardware performance monitoring capabilities, of retrieving the final performance value at block 110, setting the initial value equal to the value retrieved or reprogramming the performance monitoring unit to start counting from a predefined value if the retrieving and reprogramming procedures do not substantially affect performance monitoring results. Then, another EU needs to be selected at block 112 in order to be added to the set of EUs to accumulate data for at block 114. In case the current request is not within the set of active EUs (pending request, emulated by the EU-indicator 16), such a request may be discarded, and a zero or any predefined value may be returned as the performance monitoring result.

Thus, one embodiment of the present invention may be a system that collects performance monitoring data in one PMU counter for all execution units, and returns the performance monitoring results either each time all the EUs complete their operation, or each time a EU that happens to fall within a set of active EUs requests for completion. This means that the data collected pertains to all EUs and the total value is computed correctly (except for the described above cases of no hardware support for EU indicators) but the distribution of the final values to the EUs is considered system dependent. Still, even this implementation dependent information on the performance data distribution may be useful,

because it reflects the real-time EU interaction features and may be useful for many other types of system performance analysis.

For an exemplary embodiment of the present invention implemented in Assembler language refer to Appendix A. The Assembler code is provided for the purpose of illustration only and does not constitute a complete software performance monitoring system. Furthermore, one skilled in the art will recognize that embodiments of the present invention may be implemented in other ways and using other programming languages.

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The techniques described herein are not limited to any particular hardware or software configuration; they may find applicability in any computing or processing The techniques may be implemented in logic embodied in hardware, environment. software, or firmware components, or a combination of the above. The techniques may be implemented in programs executing on programmable machines such as mobile or stationary computers, personal digital assistants, set top boxes, cellular telephones and pagers, and other electronic devices, that each include a processor, a storage medium readable by the processor (including volatile and non-volatile memory and/or storage elements), at least one input device, and one or more output devices. Program code is applied to the data entered using the input device to perform the functions described and to generate output information. The output information may be applied to one or more output devices. One of ordinary skill in the art may appreciate that the invention can be practiced with various computer system configurations, including multiprocessor systems, minicomputers, mainframe computers, and the like. The invention can also be practiced in distributed computing environments where tasks may be performed by remote processing devices that are linked through a communications network.

Each program may be implemented in a high level procedural or object oriented programming language to communicate with a processing system. However, programs may be implemented in assembly or machine language, if desired. In any case, the language may be compiled or interpreted.

Program instructions may be used to cause a general-purpose or special-purpose processing system that is programmed with the instructions to perform the operations described herein. Alternatively, the operations may be performed by specific hardware components that contain hardwired logic for performing the operations, or by any combination of programmed computer components and custom hardware components. The methods described herein may be provided as a computer program product that may

include a machine readable medium having stored thereon instructions that may be used to program a processing system or other electronic device to perform the methods. The term "machine readable medium" used herein shall include any medium that is capable of storing or encoding a sequence of instructions for execution by the machine and that cause the machine to perform any one of the methods described herein. The term "machine readable medium" shall accordingly include, but not be limited to, solid-state memories, optical and magnetic disks, and a carrier wave that encodes a data signal. Furthermore, it is common in the art to speak of software, in one form or another (e.g., program, procedure, process, application, module, logic, and so on) as taking an action or causing a result. Such expressions are merely a shorthand way of stating the execution of the software by a processing system to cause the processor to perform an action or produce a result.

While this invention has been described with reference to illustrative embodiments, this description is not intended to be construed in a limiting sense. Various modifications of the illustrative embodiments, as well as other embodiments of the invention, which are apparent to persons skilled in the art to which the invention pertains are deemed to lie within the spirit and scope of the invention.

APPENDIX A © 2002 Intel Corporation

A code example to count the number of bus accesses from a Pentium4 processor with Hyper-Threading technology enabled.

```
;;; a function to perform arbitration
              syncHT proc near
                   ;;; IN bh == Local APIC ID
 10
                   ;;; OUT eax -> spin lock flag
                   movzx eax,bh
                   shr
                         eax, l
                         eax,[pml_sync_HT + eax]
                   lea
                        acquire_spin lock
                   call
 15
                  ret
              syncHT endp
              ;;; a function to start counting
             busproc_restart proc near
20
                  mov
                         eax, 1
                  cpuid
                  shr
                        ebx, 16
                        bl,bl ;;; no HT when zero
                  or
                  jz
                       no_HT
25
                  call
                       syncHT
                  push eax
             ;;; read ESCR
                  xor
                        eax,eax
                        edx,edx
                  xor
30
                         ecx,msr_fsb_escr0
                  mov
                 rdmsr
                 test
                       bh,01h
                 jnz
                       cpu1
```

```
;;; executing on logical CPU0
              cpu0:
              ;;; if Tx clear, program own CCCR to start counting
  5
                   ;;; eax[3..2] == T0
                   ;;; eax[1..0] == T1
                       eax,03h
                   test
                  inz
                         Tlset
                          ecx,msr fsb escr0
                   mov
 10
                   mov
                          eax,busproc_escr_mask2 OR busproc escr T0
                   wrmsr
                  ;;; clear the counter
                          eax,pml_initial_count
                  mov
                          edx,pml_initial count + 4
                  mov
 15
                         edx,0ffh
                  and
                                      ;;; 40-bit counters
                          ecx,msr_bpu counter0
                  mov
                  wrmsr
                         ecx,msr bpu cccr0
                  mov
                         eax, busproc cccr mask PMI0
                  mov
20
                  wrmsr
                  jmp
                         HT_exit
             ;;; else set T-own in ESCR
             T1set:
25
                         ecx,msr_fsb_escr0
                  mov
                        eax,busproc_escr_mask2 OR busproc escr T0
                  or
                  wrmsr
                        HT exit
                 jmp
30
             ;;; executing on logical CPU1
            cpul:
            ;;; if Tx clear, program own CCCR to start counting
                 ;;; eax[3..2] == T0
```

```
;;; eax[1..0] == T1
                   test
                        eax,0ch
                        T0set
                  jnz
                          ecx,msr_fsb_escr0
                   mov.
  5
                          eax,busproc_escr_mask2 OR busproc_escr_T1
                   mov.
                   wrmsr
                  ;;; clear the counter
                  mov
                          eax,pml_initial_count
                          edx,pml_initial_count + 4
                  mov
 10
                  and
                         edx,0ffh
                                     ;;; 40-bit counters
                          ecx,msr_bpu_counter1
                  mov
                  wrmsr
                         ecx,msr_bpu_cccrl
                  mov
                         eax,busproc_cccr_mask_PMI1
                  mov
15
                  wrmsr
                  jmp
                         HT_exit
             ;;; else set T-own in ESCR
             T0set:
20
                  mov
                         ecx,msr_fsb_escr0
                        eax,busproc_escr_mask2 OR busproc_escr_T1
                  or
                  wrmsr
             HT_exit:
                  pop
                        eax
25
                  call
                       release_spin_lock ...
                  ret
             no_HT:
                         eax,pml initial count
                  mov
                         edx,pml_initial_count + 4
                  mov
30
                 and
                        edx,0ffh
                                     ;;; 40-bit counters
                 mov
                         ecx,msr_bpu_counter0
                 wrmsr
                         ecx,msr_fsb_escr0
                 mov
```

```
mov
                         eax,busproc_escr_mask2 OR busproc_escr_T0
                  wrmsr
                         ecx,msr_bpu_cccr0
                  mov
                         eax,busproc_cccr_mask_PMI0
                  mov
 5
                  wrmsr
                  ret
             busproc_restart endp
             ;;; a function to stop counting and retrieve final value
             busproc_freeze_read proc near
10
                  ;;; OUT edx:eax = current count
                  mov
                         eax, 1
                  cpuid
                  shr
                        ebx, 16
15
                  or
                        bl,bl ;;; no HT when zero
                       no_HT
                  jz
                  call
                       syncHT
                  push eax
             ;;; read ESCR
20
                  xor
                        eax,eax
                        edx,edx
                  xor
                         ecx,msr_fsb_escr0
                  mov
                  rdmsr
                       bh,01h
                  test
25
                 jnz
                        cpu1
             ;;; executing on logical CPU0
             cpu0:
             ;;; if Tx clear, program own CCCR to stop counting
30
                  ;;; eax[3..2] == T0
                 ;;; eax[1..0] == T1
                 test
                       eax,03h
                       T1set
                 jnz
```

```
;;; stop counting
                          eax,busproc_ccr_stop_mask
                   mov.
                   xor
                         edx,edx
                          ecx,msr_bpu cccr0
                   mov
  5
                   wrmsr
                   ;;; clear ESCR
                          ecx,msr_fsb_escr0
                   mov.
                         eax,eax
                   xor
                         edx,edx
                   xor
 10
                   wrmsr
                   ;;; read count into edx:eax
                   mov
                          ecx,msr_bpu_counter0
                  rdmsr
                  jmp
                         HT_exit
 15
             ;;; else
             T1set:
             ;;; clear T-own in ESCR
                         eax, NOT busproc_escr T0
                  and
20
                  mov
                         ecx,msr fsb escr0
                  wrmsr
             ;;; read own CCCR
                  mov
                         ecx,msr_bpu_cccr0
                  rdmsr
25
                  test
                        eax,cccr_enabled
                  jz
                       disabled0
             enabled0:
                  ;;; program the other's CCCR
                  mov
                         eax,pml_initial count
30
                         edx,pml_initial_count + 4
                  mov
                  and
                        edx,0ffh
                                     ;;; 40-bit counters
                 mov
                         ecx,msr_bpu_counter1
                 wrmsr
```

```
mov
                          ecx,msr_bpu_cccr1
                          eax,busproc_ccr_mask_PMI1
                   mov
                   wrmsr
                   ;;; stop counting
  5
                          eax,busproc_ccr_stop_mask
                   mov
                         edx,edx
                   xor
                   mov
                          ecx,msr_bpu_cccr0
                   wrmsr
                   ;;; read count into edx:eax
 10
                  mov
                          ecx,msr_bpu_counter0
                  rdmsr
                  jmp
                         HT_exit
             disabled0:
 15
                  ;;; return zero count
                        edx,edx
                  xor
                  xor
                        eax,eax
                  jmp
                        HT exit
20
             ;;; executing on logical CPU1
             cpul:
             ;;; if Tx clear, program own CCCR to stop counting
                  ;;; eax[3..2] == T0
                  ;;; eax[1..0] == T1
25
                 test
                      eax,0ch
                       T0set
                 jnz
                 ;;; stop counting
                 mov
                         eax,busproc_ccr_stop_mask
                 xor
                        edx,edx
30
                 mov
                        ecx,msr_bpu_cccr1
                 wrmsr
                 ;;; clear ESCR
                 mov
                        ecx,msr_fsb_escr0
```

```
xor
                         eax,eax
                         edx,edx
                   xor
                   wrmsr
                   ;;; read count into edx:eax
  5
                   mov
                          ecx,msr_bpu_counter1
                   rdmsr
                   jmp
                         HT_exit
              ;;; else
 10
              T0set:
              ;;; clear T-own in ESCR
                         eax, NOT busproc_escr_T1
                   and
                          ecx,msr_fsb escr0
                  mov
                  wrmsr
 15
             ;;; read own CCCR
                  mov
                         ecx,msr_bpu cccrl
                  rdmsr
                  test
                        eax,cccr enabled
                       disabled1
                  jz
20
             enabled1:
                  ;;; program the other's CCCR
                  mov
                         eax,pml_initial_count
                         edx,pml_initial_count + 4
                  mov
                        edx,0ffh
                  and
                                     ;;; 40-bit counters
25
                  mov
                         ecx,msr_bpu_counter0
                  wrmsr
                  mov
                         ecx,msr_bpu cccr0
                         eax,busproc_ccr_mask_PMI0
                  mov
                  wrmsr
30
                  ;;; stop counting
                         eax,busproc_cccr_stop_mask
                  mov
                 xor
                        edx,edx
```

mov

ecx,msr_bpu_cccr1

```
wrmsr
                   ;;; read count into edx:eax
                          ecx;msr_bpu_counter1
                   rdmsr
  5
                  jmp
                         HT_exit
              disabled1:
                  ;;; return zero count
                         edx,edx
                  xor
 10
                  xor
                         eax,eax
             HT_exit:
                  xchg eax,[esp]
                  call
                        release_spin_lock
15
                  pop
                         eax
                  ret
             no_HT:
                  ;;; stop counting
20
                  mov
                         eax,busproc_cccr_stop mask
                        edx,edx
                  xor
                  mov
                         ecx,msr_bpu_cccr0
                  wrmsr
                  ;;; clear ESCR
25
                 mov
                         ecx,msr_fsb_escr0
                 xor eax, eax
                        edx,edx
                 xor
                 wrmsr
                 ;;; read count into edx:eax
30
                 mov
                         ecx,msr_bpu_counter0
                 rdmsr
                 ret
            busproc_freeze_read
                                   endp
```

CLAIMS

What is claimed is:

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1. In a system that shares performance monitoring units between multiple execution units, a method comprising:

receiving a request from an execution unit to at least one of start and stop performance monitoring operation;

maintaining a correct sequence of requests to start and stop performance monitoring operation;

performing arbitration to acquire exclusive execution for one of a plurality of request initiators;

allocating a request to start performance monitoring operation;

determining a number of allocated requests to start performance monitoring operation;

initiating performance monitoring operation; removing a request to start performance monitoring operation; determining an active mode of performance monitoring operation, and completing performance monitoring operation.

- 2. The method of claim 1, wherein maintaining the correct sequence of requests comprises ensuring that the stop request follows the start request for the same execution unit, and additional requests are ignored.
- 3. The method of claim 1, wherein allocating the request comprises indicating to the system that a request to start performance monitoring operation is pending for a specific execution unit.
- 4. The method of claim 1, wherein allocating the request and initiating of performance monitoring operation are performed if a request to start performance monitoring operation was received.
- 5. The method of claim 1, wherein removing the request comprises indicating to the system that no request to start performance monitoring operation is pending for the current execution unit.
- 6. The method of claim 1, wherein removing the request, determining active mode, and completing of performance monitoring operation are performed if a request to stop performance monitoring operation was received.

- 7. The method of claim 1, wherein initiating performance monitoring operation comprises programming the performance monitoring unit to start collecting performance data for the execution unit that requested said operation if no other request was previously allocated.
- The method of claim 7, further comprising at least one of setting an initial performance value to the current value of performance monitoring unit counter and setting the initial performance value and the performance monitoring unit counter to a predefined value.
 - 9. The method of claim 1, wherein initiating performance monitoring operation further comprises programming the performance monitoring unit to start collecting performance data for the execution unit that requested said operation, in addition to collecting performance data of other execution units which previously requested said operation if there are other requests previously allocated.

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- 10. The method of claim 1, wherein determining the active mode of performance monitoring operation comprises detecting whether programming of performance monitoring unit was performed for the execution unit that requested to stop performance monitoring operation.
- 11. The method of claim 1, wherein completing performance monitoring operation comprises retrieving final performance data and programming the performance monitoring unit to stop collecting of performance data if no other requests for the performance monitoring unit are allocated.
- 12. The method of claim 1, wherein completing performance monitoring operation further comprises programming the performance monitoring unit to stop collecting of performance data for the execution unit that requested said operation, if there are other requests previously allocated and the performance monitoring unit was in active mode for the execution unit.
- 13. The method of claim 12, further comprising retrieving current performance data, setting the initial performance value equal to the value retrieved or reprogramming the performance monitoring unit to start counting from a predefined value and setting the initial performance value equal to the predefined value.
 - 14. The method of claim 12, further comprising: selecting another execution unit;

programming the performance monitoring unit to start collecting performance data for the selected execution unit.

15. The method of claim 14, wherein selecting another execution unit comprises selecting, by external means, of a request previously allocated by another execution unit and determining the execution unit that allocated said request.

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- 16. The method of claim 14, wherein programming the performance monitoring unit further comprises enabling performance data collection for the selected execution unit in addition to performance data of other execution units which previously requested said operation.
- 17. An article comprising: a machine accessible medium having a plurality of machine readable instructions, wherein when the instructions are executed by a processor, the instructions provide for sharing of performance monitoring units between multiple execution units by:

receiving a request from an execution unit to at least one of start and stop performance monitoring operation;

maintaining a correct sequence of requests to start and stop performance monitoring operation;

performing arbitration to acquire exclusive execution for one of a plurality of request initiators;

allocating a request to start performance monitoring operation;

determining a number of allocated requests to start performance monitoring operation;

initiating performance monitoring operation; removing a request to start performance monitoring operation; determining an active mode of performance monitoring operation; and completing performance monitoring operation.

- 18. The article of claim 17, wherein instructions for maintaining the correct sequence of requests comprise instructions for ensuring that the stop request follows the start request for the same execution unit, and additional requests are ignored.
- 30 19. The article of claim 17, wherein instructions for allocating the request comprise instructions for indicating to the system that a request to start performance monitoring operation is pending for a specific execution unit.

- 20. The article of claim 17, wherein instructions for allocating the request and initiating of performance monitoring operation are executed if a request to start performance monitoring operation was received.
- 21. The article of claim 17, wherein instructions for removing the request comprise instructions for indicating to the system that no request to start performance monitoring operation is pending for the current execution unit.

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- 22. The article of claim 17, wherein instructions for removing the request, determining active mode, and completing of performance monitoring operation are executed if a request to stop performance monitoring operation was received.
- 23. The article of claim 17, wherein instructions for initiating performance monitoring operation comprise instructions for programming the performance monitoring unit to start collecting performance data for the execution unit that requested said operation if no other request was previously allocated.
- 24. The article of claim 23, further comprising instructions for at least one of setting an initial performance value to the current value of a performance monitoring unit counter and setting the initial performance value and the performance monitoring unit counter to a predefined value.
- 25. The article of claim 17, wherein instructions for initiating of performance monitoring operation further comprise instructions for programming the performance monitoring unit to start collecting performance data for the execution unit that requested said operation, in addition to performance data of other execution units which previously requested said operation if there are other requests previously allocated.
- 26. The article of claim 17, wherein instructions for determining the active mode of performance monitoring operation comprise instructions for detecting whether programming of performance monitoring unit was performed for the execution unit that requested to stop performance monitoring operation.
- 27. The article of claim 17, wherein instructions for completing performance monitoring operation comprise instructions for retrieving final performance data and programming the performance monitoring unit to stop collecting of performance data if no other requests for the performance monitoring unit are allocated.
- 28. The article of claim 17, wherein instructions for completing performance monitoring operation further comprise instructions for programming the performance monitoring unit to stop collecting of performance data for the execution unit that requested

said operation, if there are other requests previously allocated and the performance monitoring unit was in active mode for the current execution unit.

- 29. The article of claim 28, further comprising instructions for retrieving current performance data, setting the initial performance value equal to the value retrieved or reprogramming performance monitoring unit to start counting from a predefined value and setting the initial performance value equal to the predefined value.
 - 30. The article of claim 28, further comprising instructions for: selecting another execution unit;

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programming the performance monitoring unit to start collecting of performance data for the selected execution unit.

- 31. The article of claim 30, wherein instructions for selecting another execution unit comprise instructions for selecting, by external means, of a request previously allocated by another execution unit and determining the execution unit that allocated said request.
- 15 32. The article of claim 30, wherein instructions for programming the performance monitoring unit further comprise instructions for enabling performance data collection for the selected execution unit in addition to performance data of other execution units which previously requested said operation.
- 33. A system that shares performance monitoring units between multiple execution units comprising:

logic to receive a request from an execution unit to at least one of start and stop performance monitoring operation;

logic to maintain a correct sequence of requests to start or stop performance monitoring operation;

logic to perform arbitration to acquire exclusive execution for one of a plurality of request initiators;

logic to allocate a request to start performance monitoring operation;

logic to determine a number of allocated requests to start performance monitoring operation;

logic to initiate performance monitoring operation;
logic to remove a request to start performance monitoring operation;
logic to determine an active mode of performance monitoring operation; and
logic to complete performance monitoring operation.

- 34. The system of claim 33, wherein logic to maintain a correct sequence of requests comprises logic to ensure that the stop request follows the start request for the same execution unit, and additional requests are ignored.
- 35. The system of claim 33, wherein logic to allocate the request comprises logic to indicate to the system that a request to start performance monitoring operation is pending for a specific execution unit.

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- 36. The system of claim 33, wherein logic to allocate the request and initiate performance monitoring operation is activated if a request to start performance monitoring operation was received.
- 37. The system of claim 33, wherein logic to remove the request comprises logic to indicate to the system that no request to start performance monitoring operation is pending for the current execution unit.
- 38. The system of claim 33, wherein logic to remove the request, determine active mode, and complete performance monitoring operation is activated if a request to stop performance monitoring operation was received.
- 39. The system of claim 33, wherein logic to initiate performance monitoring operation comprises logic to program the performance monitoring unit to start collecting performance data for the execution unit that requested said operation if no other request was previously allocated.
- 40. The system of claim 39, further comprising logic to at least one of set an initial performance value to the current value of a performance monitoring unit counter and set the initial performance value and the performance monitoring unit counter to a predefined value.
- 41. The system of claim 33, wherein logic to initiate performance monitoring operation further comprises logic to program the performance monitoring unit to start collecting performance data for the execution unit that requested said operation, in addition to performance data of other execution units which previously requested said operation if there are other requests previously allocated.
- 42. The system of claim 33, wherein logic to determine the active mode of performance monitoring operation comprises logic to detect whether programming of performance monitoring unit was performed for the execution unit that requested to stop performance monitoring operation.

- 43. The system of claim 33, wherein logic to complete performance monitoring operation comprises logic to retrieve final performance data and to program the performance monitoring unit to stop collecting of performance data if no other requests for the performance monitoring unit are allocated.
- 44. The system of claim 33, wherein logic to complete performance monitoring operation further comprises logic to program the performance monitoring unit to stop collecting performance data for the execution unit that requested said operation, if there are other requests previously allocated and the performance monitoring unit was in the active mode for the current execution unit.
- 45. The system of claim 44, further comprising logic to retrieve current performance data, to set the initial performance value equal to the value retrieved or to reprogram performance monitoring unit to start counting from a predefined value and set the initial performance value equal to the predefined value.
 - 46. The system of claim 44, further comprising: logic to select another execution unit;

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logic to program the performance monitoring unit to start collecting performance data for the selected execution unit.

- 47. The system of claim 46, wherein logic to select another execution unit comprises logic to select, by external means, a request previously allocated by another execution unit and to determine the execution unit that allocated said request.
- 48. The system of claim 46, wherein logic to program the performance monitoring unit further comprises logic to enable performance data collection for the selected execution unit in addition to performance data of other execution units which previously requested said operation.

ABSTRACT OF THE DISCLOSURE

Efficient performance monitoring for symmetric multi-threading systems is applicable to systems that have limited performance monitoring resources and enables efficient resource sharing on a per-execution unit basis. According to embodiments of the present invention, a special indicator may be allocated for an execution unit upon its request to start performance monitoring operation. The performance monitoring unit being shared is programmed to reset its counter and to start performance monitoring operation if there is only one execution unit requesting this operation. In case there are several requests pending, an attempt is made to program the performance monitoring unit to collect performance data for a subset of execution units the hardware is capable to support.

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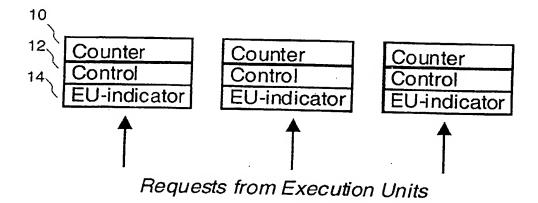
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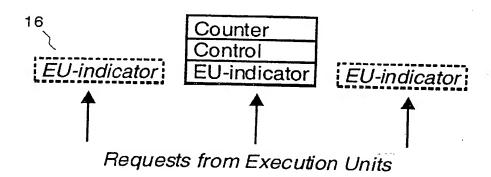
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Upon a request to stop performance monitoring operation the previously allocated indicator may be removed, and the performance monitoring unit may be programmed to stop operating if there are no more active or pending requests. The final performance data may then be obtained. Otherwise, if performance monitoring was active for the current execution unit and there are some requests pending, the performance monitoring unit may be programmed to collect data for another execution unit. In case the performance monitoring was inactive for the current execution unit, this request may be discarded, and no performance data may be returned.

Thus, the sharing of performance monitoring unit according to embodiments of the present invention results in gathering of correct information pertaining to all execution units that requested the performance monitoring operation, the information being distributed over multiple execution units on an execution time basis, that is, the first or the last execution unit that becomes inactive may take all of the data previously.





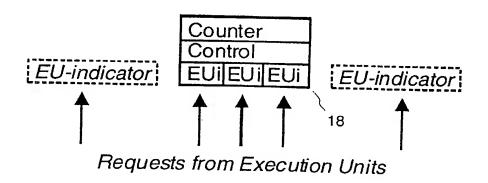


FIG. 1

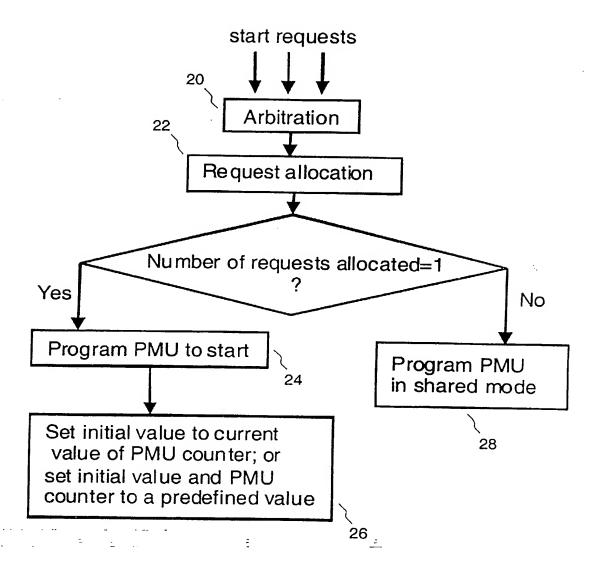


FIG. 2

